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# On Convergence of Polynomial Approximations to the Gaussian Mixture Entropy

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## Abstract

1 Gaussian mixture models (GMMs) are fundamental to machine learning due to  
2 their flexibility as approximating densities. However, uncertainty quantification  
3 of GMMs remains a challenge as differential entropy lacks a closed form. This  
4 paper explores polynomial approximations, specifically Taylor and Legendre, to  
5 the GMM entropy from a theoretical and practical perspective. We provide new  
6 analysis of a widely used approach due to Huber et al. (2008) and show that the  
7 series diverges under simple conditions. Motivated by this divergence we provide a  
8 novel Taylor series that is provably convergent to the true entropy of any GMM.  
9 We demonstrate a method for selecting a center such that the series converges from  
10 below, providing a lower bound on GMM entropy. Furthermore, we demonstrate  
11 that orthogonal polynomial series result in more accurate polynomial approxima-  
12 tions. Experimental validation supports our theoretical results while showing that  
13 our method is comparable in computation to Huber et al. We also show that in  
14 application, the use of these polynomial approximations, such as in Nonparametric  
15 Variational Inference by Gershman et al. (2012), rely on the convergence of the  
16 methods in computing accurate approximations. This work contributes useful  
17 analysis to existing methods while introducing a novel approximation supported by  
18 firm theoretical guarantees.

## 19 1 Introduction

20 Entropy is a natural measure of uncertainty and is fundamental to many information-theoretic quanti-  
21 ties such as mutual information (MI) and Kullback-Leibler (KL) divergence [8]. As a result, entropy  
22 plays a key role in many problems of ML including model interpretation [7], feature selection [6],  
23 and representation learning [27]. It is often used in the data acquisition process as in active learn-  
24 ing [25, 26], Bayesian optimal experimental design [17, 5, 3], and Bayesian optimization [13]. Yet,  
25 despite its important role entropy is difficult to calculate in general.

26 One such case is the Gaussian mixture model (GMM), where entropy lacks a closed form and is the  
27 focus of this paper. GMMs are fundamental to machine learning and statistics due to their property as  
28 universal density approximators [18]. However, the lack of a closed-form entropy requires approxima-  
29 tion, often via Monte Carlo expectation. Such stochastic estimates can be undesirable as computation  
30 becomes coupled with sample size and a deterministic approach is often preferred. Simple determinis-  
31 tic bounds can be calculated via Jensen’s inequality or Gaussian moment matching [14]. Such bounds  
32 are often too loose to be useful, leading to other options such as variational approximations [21, 10]  
33 and neural network-based approximation [2]. Yet, these deterministic estimators do not allow a  
34 straightforward tradeoff of computation and accuracy as in the Monte Carlo setting.

35 Polynomial series approximations are both deterministic and provide a mechanism for computation-  
36 accuracy tradeoff by varying the polynomial degree. In this paper we focus on three such polynomial  
37 approximations of the GMM entropy. We begin with the widely used approximation of Huber et

38 al. (2008). While this approximation yields good empirical accuracy in many settings, a proof of  
 39 convergence is lacking. In this work we show that the Huber et al. approximation in fact does  
 40 not converge in general, and we provide a divergence criterion (Theorem 3.1). In response to the  
 41 divergent behavior, we propose two alternative polynomial approximations, a Taylor and Legendre  
 42 series approximation of GMM entropy that are provably convergent. We establish in Theorem 4.2  
 43 and Theorem 4.5 that each series converges everywhere under conditions on the center point or  
 44 interval, respectively. In Theorem 4.4 we provide a simple mechanism for choosing a value to ensure  
 45 that these series converge everywhere. We additionally establish, in Theorem 4.3, that our Taylor  
 46 approximation is a convergent lower bound on the true entropy for any finite polynomial order.

47 The complexity of both Huber et al. and our proposed methods have similar computation largely  
 48 driven by polynomial order. To address this we propose an approximation that estimates the higher-  
 49 order terms by fitting a polynomial regression. This approach requires the evaluation of only three  
 50 consecutive polynomial orders to approximate higher order series. In this way we can obtain more  
 51 accurate estimates without the computational overhead of evaluating higher order polynomial terms.

52 We conclude with an empirical comparison of all polynomial approximations that produce the diver-  
 53 gent behavior of the Huber et al. approximation while our proposed methods maintain convergence.  
 54 We also compare accuracy and computation time for each method across a variety of dimensions,  
 55 number of GMM components, and polynomial orders. Finally, we show an application of our methods  
 56 in Nonparametric Variational Inference [11] where the guarantees of convergence play a large role in  
 57 the accuracy of posterior approximation via GMMs.

## 58 2 Preliminaries

59 We briefly introduce required notation and concepts, beginning with a definition of the Gaussian  
 60 mixture entropy. We will highlight the challenges that preclude efficient computation of entropy. We  
 61 conclude by defining notation that will be used for discussion of polynomial approximations.

### 62 2.1 Gaussian Mixture Entropy

63 The differential entropy of a continuous-valued random vector  $x \in \mathbb{R}^d$  with a probability density  
 64 function  $p(x)$  is given by,

$$H(p(x)) = - \int p(x) \log p(x) dx = \mathbb{E}[-\log p(x)]. \quad (1)$$

65 The differential entropy is in  $[-\infty, \infty]$  for continuous random variables. It is a measure of uncertainty  
 66 in the random variable in the sense that its minimum is achieved when there is no uncertainty in the  
 67 random vector, i.e. a Dirac delta, and approaches the maximum as the density becomes uniformly  
 68 distributed. Gaussian mixtures are ubiquitous in statistics and machine learning due to their property  
 69 as universal density approximators [18]. However, despite this flexibility, the entropy of a Gaussian  
 70 mixture requires computing the expectation of the log-sum operator, which lacks a closed form. Many  
 71 approximations and bounds are used in practice. A simple upper bound is given by the entropy of a  
 72 single Gaussian with the same mean and covariance as the mixture [14], and a lower bound can be  
 73 obtained by Jensen’s inequality. Though efficient, these bounds are very loose in practice, leading  
 74 to more useful Monte Carlo approximations, deterministic sampling [12], and numerous variational  
 75 bounds and approximations [21, 10].

### 76 2.2 Taylor Polynomials

77 In this paper we explore entropy approximation using Taylor polynomials. The  $n^{\text{th}}$ -order Taylor  
 78 polynomial of a function  $f(x)$  with evaluation point  $c$  is given by,

$$T_{f,n,c}(x) = \sum_{i=0}^n \frac{f^{(i)}(c)}{i!} (x - c)^i, \quad (2)$$

79 where  $f^{(i)}(c)$  denotes the  $i^{\text{th}}$  derivative of  $f$  evaluated at point  $c$ . The Taylor series has a region  
 80 of convergence which determines the range of  $x$ -values where the series accurately represents the  
 81 original function. It depends on the behavior of the function and its derivatives at the expansion  
 82 point. Analyzing the region of convergence is crucial for ensuring the validity of the Taylor series  
 83 approximation. Various convergence tests, such as the ratio test, help determine the  $x$ -values where  
 84 the Taylor series provides an accurate approximation.

85 **2.3 Orthogonal Polynomials**

86 Taylor series are versatile approximations, however predominately behave well near the center point  
 87 chosen. We ideally would like an approximation that performs well across a range of values. To  
 88 achieve this, we consider series approximation via orthogonal polynomials. A set of orthogonal  
 89 polynomials on the range  $[a, b]$  is an infinite sequence of polynomials  $P_0(x), P_1(x), \dots$  where  $P_n(x)$   
 90 is an  $n^{\text{th}}$  degree polynomial and for any pair of polynomials satisfies

$$\langle P_i(x), P_j(x) \rangle = \int_a^b P_i(x)P_j(x)dx = c_i\delta_{ij} \quad (3)$$

91 where  $\delta_{ij}$  is the Kronecker delta function and  $c_i$  is some constant. Orthogonal polynomials can be  
 92 used to approximate a function,  $f(x)$ , on their interval,  $[a, b]$ , by finding the projection of  $f(x)$  onto  
 93 each polynomial in the series  $P_i(x)$ .

$$f(x) = \sum_{i=1}^{\infty} \frac{\langle f(x), P_i(x) \rangle}{\langle P_i(x), P_i(x) \rangle} P_i(x) \quad (4)$$

94 Any appropriate choice of orthogonal polynomials can be used. One might be interested in consid-  
 95 ering the Chebyshev polynomials for their property of minimizing interpolation error or Legendre  
 96 polynomials for their versatility and ease of computation.

97 **3 Convergence of Polynomial Approximations**

98 To estimate the entropy  $H(p) = \mathbb{E}_p[-\log(p(x))]$  using a polynomial approximation one may ap-  
 99 proximate either the log-density  $\log(p(x))$  or just the logarithm  $\log(y)$ . We will show that estimating  
 100  $\log(p(x))$  has convergence issues and that it can be complicated to compute due to tensor arithmetic  
 101 in higher dimensions. Both of these issues will be addressed by simply approximating  $\log(y)$  and  
 102 computing the exact  $p(x)$ . All proofs are deferred to the Appendix for space.

103 **3.1 Divergence of Huber et al. Approximation**

104 We begin our exploration with a widely used approximation of the GMM entropy due to [15]. Let  $p(x)$   
 105 be a GMM and the log-GMM  $h(x) = \log(p(x))$ . Huber et al. provides a Taylor series approximation  
 106 of the GMM entropy given by,

$$\log(p(x)) = - \sum_{i=1}^M w_i \sum_{n=0}^{\infty} \frac{h^{(n)}(\mu_i)}{n!} (x - \mu_i)^n, \quad (5)$$

107 The series is  $M$  individual Taylor series evaluated at each component mean,  $\mu_i$ . The equality in  
 108 Eqn. (5) only holds if the series converges, which we will show is not the case in general.

109 **Theorem 3.1** (Divergence Criterion for Huber et al.). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x | \mu_i, \Sigma_i)$  and*  
 110 *consider the Taylor series presented in Eqn. (5). If any mean component,  $\mu_i$ , satisfies the condition*  
 111  *$p(\mu_i) < \frac{1}{2} \max(p(x))$ , then Huber et al.'s approximation diverges, otherwise it converges.*

112 Theorem 3.1 provides us with the condition that Huber et al.'s approximation Eqn. (5) will diverge.  
 113 This means that the entropy approximation will be inaccurate for any GMM with any of its modes  
 114 less than half the probability of any other point, as illustrated in Fig. 1.

115 **3.2 Taylor Series Approximation of the Logarithm**

116 Motivated by the divergence of the previous Taylor series we propose a different approach that is  
 117 provably convergent. While Huber et al. perform a Taylor decomposition of the log-GMM PDF, our  
 118 approach decomposes only the  $\log(y)$  function using a Taylor series of the function centered about  
 119 the point  $a$ . It is well-known that this series converges for values  $|y - a| < a$  and is given by,

$$\log(y) = \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^n}{na^n} (y - a)^n. \quad (6)$$

120 Note the change of  $c$  to  $a$  as the Taylor series center. This change highlights the difference in function  
 121 domains. In particular, the former series is computed on values of the random vector  $x$ , whereas ours  
 122 is computed on the PDF  $y = p(x)$ . Choosing any center  $a > \frac{1}{2} \max(p(x))$  will ensure that the series  
 123 converges everywhere.

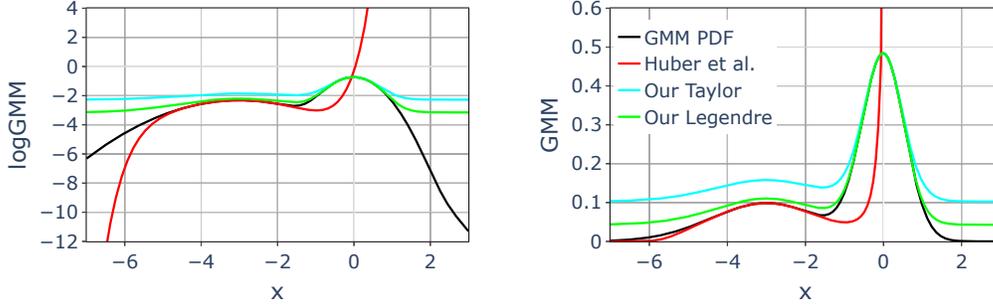


Figure 1: **Divergence of Huber et al.** and convergence of our polynomial series approximation are plotted for the Gaussian mixture,  $p(x) = .35\mathcal{N}(x | -3, 2) + .65\mathcal{N}(x | 0, .2)$ . In the left graph, the log-GMM is plotted, which is what each series is defined for. The right plot is the exponential of the series so we can see how each converge in the more familiar framework of a GMM. Notice that the Huber et al. is centered on the first component mean  $\mu_1 = -3$  and diverges around the mean of the second component  $\mu_2 = 0$  as supported by Theorem 3.1 since the mode  $\mu_1$  is less than half the probability at the mode  $\mu_2$ . Both of our methods are convergent, the Taylor series is a bound (Theorem 4.3) while the Legendre series has a lower global error.

124 **Lemma 3.2** (Convergent Taylor Series of Log). *If  $a > \frac{1}{2}\max(p(x))$ , then for all  $x$*

$$\log(p(x)) = \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{na^n} \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} p(x)^k \quad (7)$$

125 The proof of Lemma 3.2 is a simple ratio test. The only assumption on  $p(x)$  is that it has a finite  
 126 maximum, which is true for any non-degenerate GMM with positive definite component covariances.  
 127 As a result, the Taylor series converges for all  $x$  regardless of the GMM form.

### 128 3.3 Legendre Series Approximation of the Logarithm

129 For the orthogonal polynomial approximation, we consider the Legendre polynomials, specifically  
 130 the shifted Legendre polynomials [4] which are orthogonal on  $[0, a]$ ,

$$P_n(y) = L_{[0,a],n}(y) = \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} y^k \quad (8)$$

131 **Lemma 3.3** (Convergent Legendre Series of Log). *If  $a > \max(p(x))$ , and consider the shifted  
 132 Legendre polynomials on the interval  $[0, a]$  in Eqn. (8). Then for all  $x$*

$$\log(p(x)) = \sum_{n=0}^{\infty} (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! ((j+1)!)^2} L_{[0,a],n}(p(x)) \quad (9)$$

133 Again, all that is assumed about this approximation is that the max of a GMM can be bounded, so  
 134 this approximation converges for all GMMs regardless of structure.

## 135 4 GMM Entropy Approximations

136 Having established multiple polynomial approximations in Sec. 3, we now consider applying them to  
 137 the definition of entropy for a GMM. We can directly substitute the series approximation into the  
 138 entropy definition,  $H(p(x)) = \mathbb{E}_p[-\log(p(x))]$ , and push the expectation through the summations.

### 139 4.1 Huber et al. Entropy approximation

140 Applying Huber et al.'s Taylor series approximation of the  $\log(p(x))$ , we see the GMM entropy can  
 141 be approximated by,

$$H(p(x)) = - \sum_{i=1}^m w_i \sum_{n=0}^{\infty} \frac{h^{(n)}(\mu_i)}{n!} \mathbb{E}_{q_i} [(x - \mu_i)^n], \quad (10)$$

142 where  $q_i(x) = \mathcal{N}(x | \mu_i, \Sigma_i)$  is shorthand for the  $i^{\text{th}}$  Gaussian component. The attractive feature  
 143 of Eqn. (10) is that it simplifies the expected value of a log-GMM to the  $n^{\text{th}}$  central moments of

144 the  $i^{\text{th}}$  component which, is exactly zero when  $n$  is odd and has a closed form when  $n$  is even.  
 145 However, this approximation had some major limitations. Theorem 3.1 shows that this approximation  
 146 is not guaranteed to converge which is supported by experimental results in Sec. 6. Furthermore, in  
 147 higher dimensions,  $h^{(n)}(\mu_i) = \frac{\partial^n h(\mu_i)}{\partial x_1^{j_1} \dots \partial x_d^{j_d}}$ , where  $j_1 + \dots + j_d = n$  which grows rapidly, is an  $n$   
 148 dimensional tensor. This is cumbersome to compute and is difficult to deal with the tensor arithmetic  
 149 required beyond a Hessian. In practice, this limits Eqn. (10) to only second order approximations (or  
 150 third order since the third centered moment is zero) unless we are in one dimension.

#### 151 4.2 Taylor Series Entropy Approximation

152 Having established the convergent Taylor series of the logarithm in Lemma 3.2, we can applying the  
 153 approximation and push the expectation through the summations. This reduces the computation of the  
 154 entropy to computing  $\mathbb{E}_p[p(x)^k]$  for all  $k < n$  where  $n$  is the order of the polynomial approximation.

155 **Lemma 4.1** (Closed form expectation of powers of GMMs). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a  
 156 GMM and  $k$  be a non-negative integer. Then*

$$\mathbb{E}_p[p(x)^k] = \sum_{j_1 + \dots + j_M = k} \binom{k}{j_1, \dots, j_M} \sum_{i=1}^M w_i \left( \frac{\mathcal{N}(0|\mu_i, \Sigma_i)}{\mathcal{N}(0|\mu, \Sigma)} \prod_{t=1}^M (w_t \mathcal{N}(0|\mu_t, \Sigma_t)^{j_t}) \right) \quad (11)$$

157 where  $\Sigma = (\Sigma_i^{-1} + \sum_{t=1}^M j_t \Sigma_t^{-1})^{-1}$  and  $\mu = \Sigma(\sum_{i=1}^M \mu_i + \sum_{t=1}^M j_t \Sigma_t^{-1} \mu_t)$ .

158 While Eqn. (11) may seem complicated at first glance, it is straightforward to compute. All terms  
 159 are Gaussian densities, polynomial functions, and binomial coefficients. Lemma 4.1 is defined for  
 160  $\mathbb{E}_p[p(x)^k]$  but an analogous definition holds for  $\mathbb{E}_p[q(x)^k]$  allowing us to apply all the following  
 161 results not only to entropy, but cross-entropy, KL and MI of GMMs. This is not the focus of this  
 162 paper, however a discussion can be found in A.4 for completeness. Using Lemma 3.2 and Eqn. (11),  
 163 we can obtain the following approximation,

$$\hat{H}_{N,a}^T(p(x)) = -\log(a) - \sum_{n=1}^N \frac{(-1)^{n-1}}{na^n} \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} \mathbb{E}_p[p(x)^k] \quad (12)$$

164 To ensure the expected value can be pushed through the infinite sum of the series, we check that our  
 165 finite order entropy approximation does still converge to the true entropy.

166 **Theorem 4.2** (Convergence of  $\hat{H}_{N,a}^T(p(x))$ ). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a GMM and  
 167 choose a Taylor center such that  $a > \frac{1}{2} \max(p(x))$ . Then, for  $\hat{H}_{N,a}^T(p(x))$  defined in Eqn. (12)*

$$\lim_{N \rightarrow \infty} \hat{H}_{N,a}^T(p(x)) = H(p(x)) \quad (13)$$

168 Having established convergence of our estimator, it remains to provide a method for selecting a  
 169 Taylor center that meets the convergence criterion  $a > \frac{1}{2} \max(p(x))$ . In fact, we show in Theorem 4.3  
 170 that selecting a looser condition  $a > \max(p(x))$  ensures convergence from below, thus yielding a  
 171 lower bound on the true entropy.

172 **Theorem 4.3** (Taylor Series is Lower Bound of Entropy). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  and  
 173  $a > \max(p(x))$ . Then, for all finite  $N$ ,*

$$\hat{H}_{N,a}^T(p(x)) \leq H(p(x)) \quad (14)$$

174 We have now established that the Taylor center chosen as  $a > \max(p(x))$  is both convergent and  
 175 yields a lower bound. In fact, it is easy to find such a point by upper bounding the maximum of a  
 176 GMM as given in Theorem 4.4.

177 **Theorem 4.4** (Upper bound on maximum of a GMM). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$ , then*

$$\max(p(x)) \leq a = \sum_i^M w_i |2\pi\Sigma_i|^{-\frac{1}{2}} \quad (15)$$

178 In our experience choosing a center closer to the convergence criterion  $a > \frac{1}{2} \max(p(x))$  yields  
 179 slightly more accurate estimates, but not significantly so. We find in practice that performing an  
 180 approximation of the limit of the Taylor polynomial (similar to Richardson extrapolation [22]) is a  
 181 better approach for higher accuracy.

$$\hat{H}_{N,a}^{TL}(p(x)) = \hat{H}_{N-2,a}^T(p(x)) - \frac{(\hat{H}_{N-1,a}^T(p(x)) - \hat{H}_{N-2,a}^T(p(x)))^2}{\hat{H}_{N,a}^T(p(x)) - 2\hat{H}_{N-1,a}^T(p(x)) + \hat{H}_{N-2,a}^T(p(x))} \quad (16)$$

182 This limit requires three consecutive terms of the Taylor series to be computed and assumes that the  
 183 series converges at the rate  $\beta \cdot \alpha^n + \eta$  for some  $\beta < 0$  and  $0 < \alpha < 1$ . This does not hold true in  
 184 general but in practice provides higher accuracy. Derivation and discussion can be found in A.5.

185 **4.3 Legendre Entropy Approximation**

186 Now, starting with the convergent Legendre approximation considered in Lemma 3.3 and Eqn. (11),  
 187 we can obtain the following approximation,

$$\hat{H}_{N,a}^L(p(x)) = - \sum_{n=0}^N (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! (j+1)!^2} L_{[0,a],n}(\mathbb{E}_p[p(x)^k]) \quad (17)$$

188 Again, we check that taking the expectation of our series does not effect convergence.

189 **Theorem 4.5** (Convergence of  $\hat{H}_{N,a}^L(p(x))$ ). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a GMM and  
 190 choose an interval such that  $a > \max(p(x))$ . Then for  $\hat{H}_{N,a}^L(p(x))$  defined in Eqn. (17)*

$$\lim_{N \rightarrow \infty} \hat{H}_{N,a}^L(p(x)) = H(p(x)) \quad (18)$$

191 Now having established converge criterion for the Legendre series approximation, we need to choose  
 192 an upper point of the interval for the Legendre series. We need to choose  $a > \max(p(x))$  which is  
 193 satisfied by the same  $a$  found in Lemma 4.4.

194 **5 Related Work**

195 Numerous approximation methods exist in the literature for estimating entropy and related information  
 196 measures, such as mutual information and Kullback-Leibler divergence, in the context of Gaussian  
 197 Mixture Models (GMMs). Monte Carlo estimation, deterministic bounds using Jensen’s inequality,  
 198 best-fit moment matched Gaussians, and numerical integral approximations based on the uncentred  
 199 transform have been explored [16, 14, 15]. This paper focuses on the Taylor approximation by Huber  
 200 et al., an alternative Taylor approximation is proposed by Sebastiani [24], which assumes a shared  
 201 covariance matrix among GMM components in the high variance setting. However, neither Huber  
 202 et al. or Sebastiani provide theoretical analysis or convergence guarantees offered in our present  
 203 work. An analysis conducted by Ru et al. [23] explores the efficiency of Huber et al.’s method and  
 204 demonstrates that deterministic quadrature methods can be equally fast and accurate in a single  
 205 dimension, however quadrature methods scale poorly with dimension, at  $\mathcal{O}(N^D)$  where  $N$  is the  
 206 number of quadrature points per dimension and  $D$  is the dimension of the problem.

207 Variational approximations and bounds are also widely explored for estimating entropy and mutual  
 208 information (MI). Much of this work is motivated by the use of Gibbs’ inequality, which leads to  
 209 bounds on entropy and MI [1]. Later work explored similar techniques for upper and lower bounds  
 210 on MI [21, 10]. More recent work uses artificial neural networks (ANNs) as function approximators  
 211 for a variety of information-theoretic measures based on differential entropy. The MI neural estimator  
 212 (MINE) uses such an approach for representation learning via the *information bottleneck* [2] based on  
 213 the Donsker-Varadhan (DV) lower bound on KL [9]. Related methods use ANNs for optimizing the  
 214 convex conjugate representation of Nguyen et al. [20]. McAllester and Stratos [19] show that many of  
 215 these distribution-free approaches based on ANN approximation rely on Monte Carlo approximations  
 216 that have poor bias-variance characteristics which they provide their own Difference of Entropies  
 217 (DoE) estimator that achieves the theoretical limit on estimator confidence.

218 **6 Experiments**

219 We consider two experiments, a synthetic GMM section where we look at divergence of Huber et  
 220 al. approximation (Eqn. (10)) and convergence of our three methods, our Taylor (Eqn. (12)), Taylor  
 221 limit (Eqn. (16)), and our Legendre (Eqn. (17)). Furthermore, we give comparisons of accuracy and  
 222 computation time across a variety of setting of approximation order, number of GMM components,  
 223 and dimension for all methods. We then show our how our methods can be applied in practice to  
 224 Nonparametric Variational Inference [11] where the convergence guarantees of the estimators has a  
 225 noticeable accuracy improvement on their algorithm.

226 **6.1 Synthetic Multivariate GMM**

227 To highlight the theoretical properties, such as convergence, divergence, accuracy, and lower-bound  
 228 of methods as discussed in Sec. 4, we will consider some synthetic GMMs. We create two GMMs  
 229 similar to the example published in [15] (original experiment recreated in A.6). We consider a single  
 230 and multi-dimensional dimensional case that satisfy the divergence criterion in Theorem 3.1. We also  
 231 look at a time and accuracy analysis versus dimension, components, and polynomial order.

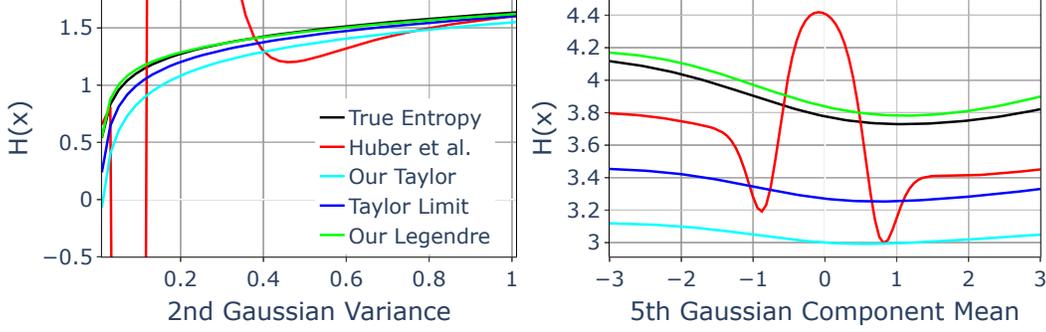


Figure 2: **Scalar GMM** example is plotted on the left. The variance of a component of a two component GMM is varied between  $\sigma_2^2 \in (0, 1]$  as in theory according to Theorem 3.1, the example will be divergent where  $\sigma_2^2 < .46$  and convergent above. We plot the fourth order of each method and see that Huber et al.’s approximation does diverge where the theory predicts. **Two dimensional GMM** with five components is consider on the right. Here the mean of a single component is shifted from  $\mu_5 = [-3, -3]^T$  to  $\mu_5 = [3, 3]^T$ . We consider the third order approximation of each method and see that Huber et al. is poorly behaved. In both examples, we see that our Taylor method is a lower bound (Theorem 4.3), the Taylor limit provides higher accuracy, and Our Legendre method is a highly accurate approximate.

232 **Scalar GMM** In this experiment, we consider a scalar GMM as fourth order and above cannot  
 233 be easily computed in higher dimensions for Huber et al. due to tensor arithmetic. We use a simple  
 234 two-component GMM with parameters  $w_1 = 0.35$ ,  $w_2 = 0.65$ ,  $\mu_1 = -2$ ,  $\mu_2 = -1$ ,  $\sigma_1^2 = 2$ , and  
 235  $\sigma_2^2 \in (0, 1]$ . We are changing the variance of the second Gaussian,  $\sigma_2^2$ , in the range  $(0, 1]$  because  
 236 the condition for divergence in Theorem 3.1 ( $p(x = \mu_1) < \frac{1}{2}p(x = \mu_2)$ ) is satisfied approximately  
 237 when  $\sigma_2^2 < 0.46$  meaning this experiment should have regions of both convergence and divergence  
 238 for Huber et al. approximation. Fig. 2 (left) shows the fourth order approximations of all methods.  
 239 We see that the Huber et al. approximations diverges as expected in the range where  $\sigma_2^2 < .46$ . Our  
 240 Taylor method remains convergent and accurate for all values while maintaining a lower bound.  
 241 Again, our limit method gains us some accuracy and still manages to be lower bound. In this case,  
 242 we see that the Legendre approximation is a near perfect fit for the entropy.

243 **Multivariate GMM** To demonstrate that divergence is not limited to single dimension or higher  
 244 orders, we consider a five-component, two-dimensional GMM with the parameters  $w_i = .2 \forall i$ ,  
 245  $\mu_1 = [0, 0]^T$ ,  $\mu_2 = [3, 2]^T$ ,  $\mu_3 = [1, -.5]^T$ ,  $\mu_4 = [2.5, 1.5]^T$ ,  $\mu_5 = c[1, 1]^T$  for  $c \in [-3, 3]$ ,  
 246  $\Sigma_1 = .25\mathbf{I}_2$ ,  $\Sigma_2 = 3\mathbf{I}_2$ , and  $\Sigma_3 = \Sigma_4 = \Sigma_5 = 2\mathbf{I}_2$  where  $\mathbf{I}_2$  is the two dimensional identity  
 247 matrix. This examples shifts the mean of the fifth component to show that simply the location of  
 248 components can make the Huber et al. approximation behave poorly. Fig. 2 (right) shows the third  
 249 order approximation of each method. We see that Huber et al. is clearly not well behaved in this case  
 250 even with low order approximation. Furthermore, we continue to see a lower bound by our Taylor  
 251 method, an increased accuracy from out limit method, and that the Legendre approximation is very  
 252 close to the true entropy.

### 253 6.1.1 Computation Time

254 In this experiment we empirically analyze the computation time of each method as a function of  
 255 Gaussian dimension, number of Gaussian components, and the order of each polynomial approximation.  
 256 The baseline of each method will be compared to the Monte Carlo estimation of entropy using  
 257  $L = 1000$  samples  $\{x_j\}_{j=1}^L \sim p$ . The Monte Carlo estimator is given by  $\hat{H} = \frac{1}{L} \sum_j (-\log p(x_j))$ .

258 **Dimension** In Fig. 3 (left), we evaluate the accuracy and computation time for 30 two-component  
 259 GMMs per dimension in the range of  $[1, 50]$ . Comparing second order approximations of all methods  
 260 against the Monte Carlo estimator, our polynomial approximations demonstrate similar accuracy and  
 261 nearly identical computation time. The results are comparable to Huber, indicating that our methods  
 262 preserve accuracy and computation efficiency while providing convergence guarantees.

263 **GMM Components** In Fig. 3 (middle), accuracy and computation time are presented for 30  
 264 two-dimensional GMMs with varying numbers of components (from 1 to 20) using second order  
 265 approximations. Legendre and Huber methods show slightly higher accuracy compared to our  
 266 Taylor approximation and Taylor limit. Notice, Huber’s standard deviation also increases with more  
 267 components, due to the increased likelihood of satisfying the divergence condition in Theorem 3.1.  
 268 Computation time remains similar for all methods, but is more prohibitive for higher components.

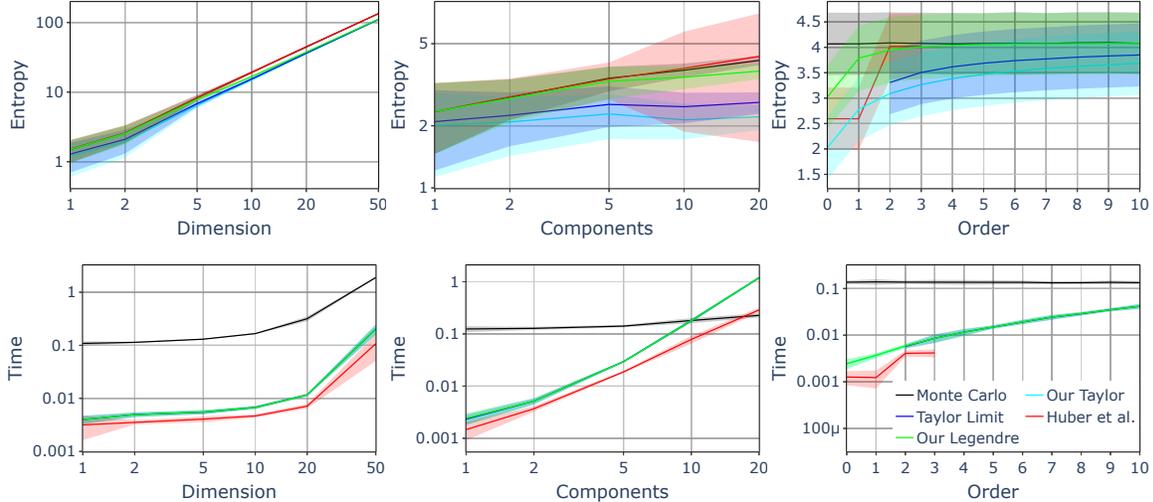


Figure 3: **Dimension (left)** of a two-component GMM varies from zero to fifty, for the second order of each method. Our methods show comparable accuracy and computation time to Huber, regardless of dimension. **Number of components (middle)** in a two-dimensional GMM is considered for the second order approximation of all methods. Huber et al. and our Legendre approximations are nearly equivalent in accuracy, while the Taylor series and Taylor limit serve as lower bounds. Computation time for all our methods is identical and comparable to Huber et al., deviating only at high numbers of components. **Order (right)** of each approximation is varied for a three-dimensional, two-component GMM. Huber et al. is plotted up to order three, as higher orders are restrictive due to tensor arithmetic and Taylor limit starts at order two as it requires three consecutive terms.

269 **Polynomial Order** Fig. 3 (right) shows as the order of the polynomial approximation increases for  
 270 two-component GMMs in three dimensions. Legendre and Huber methods show higher accuracy  
 271 compared to Taylor approximation and Taylor limit. Huber is limited to order 3 due to Tensor  
 272 arithmetic, while Taylor limit starts at order 2 as it requires multiple orders. Computation times are  
 273 similar across all methods. Notice no accuracy is gained from zero to first order and from second to  
 274 third order in Huber’s approximation due to relying on odd moments of Gaussians which are zero.

## 275 6.2 Nonparametric Variational Inference

276 Consider a target density  $p(x, \mathcal{D})$  with latent variables  $x$  and observations  $\mathcal{D}$ . The NPV approach [11]  
 277 optimizes the evidence lower bound (ELBO),  $\log p(x, \mathcal{D}) \geq \max_q H_q(p(x, \mathcal{D})) - H_q(q(x)) \equiv \mathcal{L}(q)$   
 278 w.r.t. an  $m$ -component GMM variational distribution  $q(x) = \frac{1}{N} \sum_{i=1}^m \mathcal{N}(x|\mu_i, \sigma_i^2 I_d)$ . The GMM  
 279 entropy lacks a closed-form so NPV applies Jensen’s lower bound as an approximation,  $\hat{H}_q^J(q(x))$ .  
 280 The cross entropy also lacks a closed-form, so NPV approximates this term using the analogous  
 281 Huber et al. Taylor approximation. Specifically, NPV expands the log density around the means of  
 282 each GMM component as,

$$H_q(p(x)) \approx - \sum_{i=1}^M w_i \sum_{n=0}^N \frac{\nabla^2 \log(p(\mu_i))}{n!} \mathbb{E}_{q_i} [(x - \mu_i)^n] = \hat{H}_{N,q}^H(p(x)) \quad (19)$$

283 However, Eqn. (19) is subject to the divergence criterion of Theorem 3.1 if  $2p(\mu_i) \leq \max(p(x))$ .  
 284 By replacing the entropy terms with our convergent series approximations we observe significant  
 285 improvements in accuracy.

286 **In our approach**, we will highlight and address two problems with the NPV algorithm; the potential  
 287 divergence of  $\hat{H}_{N,q}^H(p(x))$  and the poor estimation of the GMM entropy via  $\hat{H}_q^J(q(x))$ . To address  
 288 the potential divergence of  $\hat{H}_{N,q}^H(p(x))$ , we will take motivation from the results found in [23] and  
 289 use a 2 point Gauss-Hermite quadrature method to approximate  $H_q(p(x))$ . This method will be  
 290 a limiting factor in scaling the NPV algorithm in dimension, however it guarantees that the cross-  
 291 entropy approximation will not diverge. This alteration leads to a solution for the inconsistency of  
 292 the ELBO approximations. Then, Jensen’s inequality is a very poor approximation for entropy in  
 293 general, instead we will use the three methods we have introduced, our Taylor, Taylor limit, and our  
 294 Legendre, as the GMM entropy approximations for higher accuracy. Fig. 4 shows an approximation  
 295 of a two dimensional, three component mixture Student T distribution using a five component GMM  
 296 in the traditional NPV, our modified NPV algorithm with our Taylor and Legendre approximation.

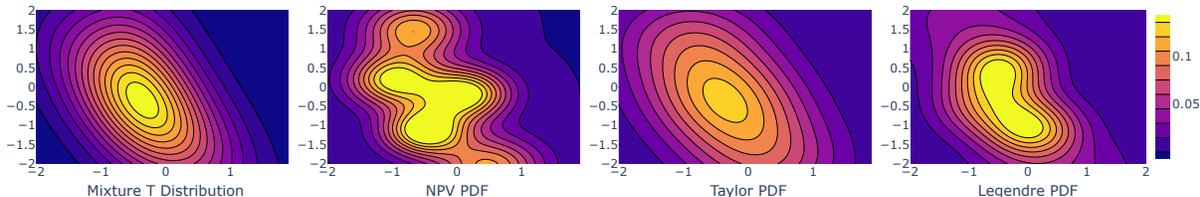


Figure 4: A three component mixture Student T distribution PDF (far-left) is approximated by a five component GMM using traditional NPV (left), our algorithm using a 6<sup>th</sup> order Taylor polynomial (right), and Legendre polynomial (far-right). We see that NPV both has issues with finding correct placement of means and sets the variances of the GMM components to be too narrow. Our methods do a better job of assigning means and the Legendre method seems to set the variances slightly better than our Taylor.

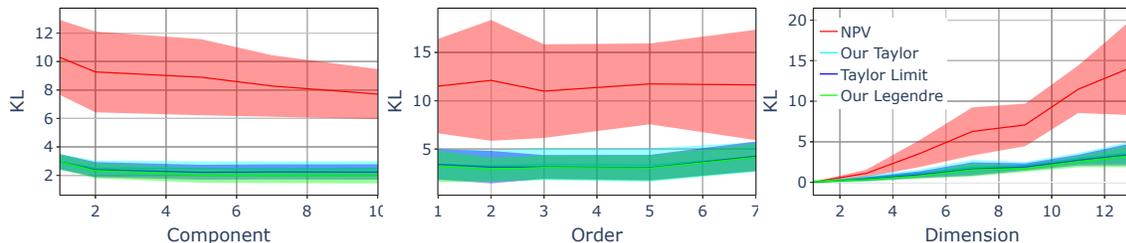


Figure 5: The above figures show the accuracy of each method across varying components, orders, and dimensions in approximating a multivariate mixture T distribution with a GMM. Our method consistently improves accuracy significantly. Low order of the convergent estimators provide substantial approximation improvement (middle). Most accuracy improvements are achieved with a small number of components, unlike NPV (left) which continues to need higher number of components to see good accuracy return. The guaranteed convergence of the approximation in higher dimensions seems to have a drastic improvement on accuracy ().

297 **The results**, as seen in Fig. 5, highlight the accuracy of each method versus the number of components,  
 298 the order of our polynomial approximation, and the dimension of the GMM. In each experiment,  
 299 we are approximating a multivariate mixture T distribution,  $p(x)$ . We randomize the parameters  
 300 of  $p(x)$  and the initialization parameters of the variational GMM,  $q(x)$ , for optimization. The KL  
 301 is approximated using a 100000 Monte Carlo approximation after convergence of each algorithm.  
 302 We see that in all cases of components, order, and dimension, our method achieves significant  
 303 accuracy improvements. We see that we can use low order approximations to receive substantial  
 304 approximation improvement (Fig. 5 (middle)). We see all methods gain accuracy as number of  
 305 components increase (Fig. 5 (left)) however our methods see most of the accuracy improvements with  
 306 only a few components, whereas NPV has substantially worse approximations with low components.  
 307 Finally, we see we maintain a lower variance and KL than NPV with all our methods as the dimension  
 308 grows (Fig. 5 (right)). For further discussion of the experiment, see A.7.

## 309 7 Discussion

310 We have provided novel theoretical analysis of the convergence for the widely used Huber et al. Taylor  
 311 approximation of GMM entropy and established that the series diverges under conditions on the  
 312 component means. We address this divergence by introducing multiple novel methods which provably  
 313 converge. We wish to emphasize that the Huber et al. approximation tends to yield accurate results  
 314 when it is convergent and the intention of this work is not to dissuade the use of this approximator.  
 315 Quite the contrary, this work encourages the use of either Huber et al. or our own estimator by  
 316 providing a solid theoretical foundation for both methods. We acknowledge that there are contexts in  
 317 which one method may be preferred over the other, for example when bounds are preferred, or when  
 318 convergence criteria are provably satisfied.

319 There are several areas that require further investigation. For example, one limitation of both methods  
 320 is that they scale poorly with polynomial order and number of components. In fact, Huber et al. cannot  
 321 easily be calculated for fourth order and above, due to tensor arithmetic. Our approximation works  
 322 well in practice, but is limited solely to GMM densities. Further work is necessary to efficiently apply  
 323 our convergent series to situations of cross-entropy's that contain non-GMM distributions.

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398 **A Appendix**

399 **A.1 Section 3 Proofs**

400 **Theorem 3.1** (Divergence Criterion for Huber et al.). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x | \mu_i, \Sigma_i)$  be a GMM*  
 401 *and consider the Taylor series presented by Huber et al. in Eqn. (5). If any mean component,  $\mu_i$ ,*  
 402 *satisfies the condition  $p(\mu_i) < \frac{1}{2} \max(p(x))$ , then Huber et al.'s approximation diverges, otherwise it*  
 403 *converges.*

404 *Proof.* Let  $f(y) = \log(y)$ ,  $g(x) = p(x)$ , and  $h(x) = f(g(x)) = \log(p(x))$ . Huber et al. creates the  
 405 Taylor series in Eqn. (5) with  $N^{\text{th}}$  order approximation

$$\sum_{i=1}^M w_i T_{h,N,\mu_i}(x) = \sum_{i=1}^M w_i \sum_{n=0}^N \frac{h^{(n)}(\mu_i)}{n!} (x - \mu_i)^n, \quad (20)$$

406 Let us consider just a single one of the Taylor series in Eqn. (20)

$$T_{h,N,\mu_i}(x) = \sum_{n=0}^N \frac{h^{(n)}(\mu_i)}{n!} (x - \mu_i)^n, \quad (21)$$

407 By Theorem 3.4 in Lang<sup>1</sup>, the Taylor series of a composition of function is equivalent to the  
 408 composition of each components Taylor series, i.e.  $T_{h,N,\mu_i}(x) = T_{f,N,g(\mu_i)} \circ T_{g,N,\mu_i}(x)$  where  $\circ$  is  
 409 the composition operation. Since  $p(x)$  is a GMM, is the sum of entire functions, and thus itself is  
 410 entire, meaning it's Taylor series,  $T_{g,N,\mu_i}(x)$ , converges everywhere. We turn our attention to the  
 411 Taylor approximation of  $\log$ ,  $T_{f,N,g(\mu_i)}$ ,

$$T_{f,N,p(\mu_i)}(x) = \log(p(\mu_i)) + \sum_{n=1}^N \frac{(-1)^{n-1}}{np(\mu_i)^n} (y - p(\mu_i))^n, \quad (22)$$

412 We can look at the  $n^{\text{th}}$  term in the series,  $b_n = \frac{(-1)^{n-1}}{np(\mu_i)^n} (y - p(\mu_i))^n$ , and use the ratio test to define  
 413 convergence.

$$\begin{aligned} \lim_{n \rightarrow \infty} \left| \frac{b_{n+1}}{b_n} \right| &= \lim_{n \rightarrow \infty} \left| \frac{(-1)^n}{(n+1)p(\mu_i)^{n+1}} (y - p(\mu_i))^{n+1} \frac{np(\mu_i)^n}{(-1)^{n-1}} (y - p(\mu_i))^{-n} \right| \\ &= \lim_{n \rightarrow \infty} \left| \frac{(-1)^n}{(-1)^{n-1}} \frac{np(\mu_i)^n}{(n+1)p(\mu_i)^{n+1}} \frac{(y - p(\mu_i))^{n+1}}{(y - p(\mu_i))^n} \right| \\ &= \lim_{n \rightarrow \infty} \left| \frac{n}{(n+1)p(\mu_i)} (y - p(\mu_i)) \right| \\ &= \lim_{n \rightarrow \infty} \frac{n}{(n+1)} \left| \frac{y - p(\mu_i)}{p(\mu_i)} \right| = \left| \frac{y - p(\mu_i)}{p(\mu_i)} \right| = L \end{aligned}$$

414 The ratio test states that the series converges if the limit,  $L$ , is strictly less than 1. However, setting  
 415  $L = 1$  and some simple manipulation, we find

$$\left| \frac{y - p(\mu_i)}{p(\mu_i)} \right| < 1 \quad \Rightarrow \quad |y - p(\mu_i)| < p(\mu_i) \quad \Rightarrow \quad y < 2p(\mu_i) \quad (23)$$

416 This only converges if all  $y < 2(p(\mu_i))$ . Consider the maximum,  $\max(p(x))$ , if it satisfies this  
 417 condition, so will every other point, if it doesn't satisfy this point, then the series is divergent by the  
 418 ratio test. So we have the convergent criterion  $\max(p(x)) < 2p(\mu_i)$ , or written in terms of divergence  
 419 criterion, we diverge if  $p(\mu_i) < \frac{1}{2} \max(p(x))$ .  $\square$

420 **Lemma 3.2** (Convergent Taylor Series of Log). *If  $a > \frac{1}{2} \max(p(x))$ , then for all  $x$*

$$\log(p(x)) = \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{na^n} \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} p(x)^k \quad (24)$$

<sup>1</sup>Lang, Serge. Complex Analysis. 4th ed. Springer, 2013. ISBN 978-1-4757-3083-8.

421 *Proof.* Consider the  $n^{\text{th}}$  term in the sum

$$b_n = \frac{(-1)^{n-1}}{na^n} (p(x) - a)^n$$

422 The ratio test says if the limit of the absolute value of successive terms converges to a value strictly  
423 less than one, then the series converges

$$\begin{aligned} \lim_{n \rightarrow \infty} \left| \frac{b_{n+1}}{b_n} \right| &= \lim_{n \rightarrow \infty} \left| \frac{(-1)^n}{(n+1)a^{n+1}} (p(x) - a)^{n+1} \frac{na^n}{(-1)^{n-1}} (p(x) - a)^{-n} \right| \\ &= \lim_{n \rightarrow \infty} \left| \frac{(-1)^n}{(-1)^{n-1}} \frac{na^n}{(n+1)a^{n+1}} \frac{(p(x) - a)^{n+1}}{(p(x) - a)^n} \right| \\ &= \lim_{n \rightarrow \infty} \left| \frac{n}{(n+1)a} (p(x) - a) \right| \\ &= \lim_{n \rightarrow \infty} \frac{n}{(n+1)} \left| \frac{p(x) - a}{a} \right| = \left| \frac{p(x) - a}{a} \right| = L \end{aligned}$$

424 We see that  $L < 1 \forall x$  iff  $a > \frac{1}{2} \max(p(x))$  in which case the series converges everywhere.  $\square$

425 **Lemma 3.3** (Convergent Legendre Series of Log). *If  $a > \max(p(x))$ , and consider the  $n^{\text{th}}$  shifted*  
426 *Legendre polynomial on the interval  $[0, a]$  in Eqn. (8). Then for all  $x$*

$$\log(p(x)) = \sum_{n=0}^{\infty} (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! ((j+1)!)^2} L_{[0,a],n}(p(x)) \quad (25)$$

427 *Proof.* Orthogonal polynomials can approximate any function that is continuous and square-integrable  
428 (see Trefethen and Bau<sup>2</sup>). In our case,  $L_{[0,a],n}(y)$  live on  $L_2([0, a])$  (referring to the second Lebesgue  
429 space on the interval  $[0, a]$ ). This means all we have to show is that  $\log(y)$  lives in this domain which  
430 means  $\|\log(y)\|_2^2 < \infty$

$$\|\log(y)\|_2^2 = \int_0^a \log(y)^2 dy = a((\log(a) - 2) \log(a) + 2) < \infty \quad (26)$$

431 So we see that  $\log(y) \in L_2([0, a])$  and therefore it's Legendre series is convergent.

432 For completeness, we now derive the Legendre series for  $\log(y)$ . We will appeal to Eqn. (8),

433  $L_{[0,a],n} = \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} y^k$ , and  $\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle = \frac{a}{2n+1}$  as found in [4]

$$\begin{aligned} \log(p(x)) &= \sum_{n=0}^{\infty} \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) \\ &= \sum_{n=0}^{\infty} \frac{2n+1}{a} \int_0^a \log(y) \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} y^k dy L_{[0,a],n}(p(x)) \\ &= \sum_{n=0}^{\infty} \frac{2n+1}{a} \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} \int_0^a \log(y) y^k dy L_{[0,a],n}(p(x)) \\ &= \sum_{n=0}^{\infty} \frac{2n+1}{a} \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)! a^{k+1} ((k+1) \log(a) - 1)}{(n-k)! (k!)^2 a^k (k+1)^2} L_{[0,a],n}(p(x)) \\ &= \sum_{n=0}^{\infty} 2n+1 \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)! ((k+1) \log(a) - 1)}{(n-k)! ((k+1)!)^2} L_{[0,a],n}(p(x)) \end{aligned}$$

434 which we know is convergent from the above discussion  $\square$

<sup>2</sup>Trefethen, Lloyd N., and David Bau III. Numerical Linear Algebra. SIAM, 1997.

435 **A.2 Section 4 Proofs**

436 **Lemma 4.1** (Closed form expectation of powers of GMMs). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a*  
 437 *GMM and  $k$  be a non-negative integer. Then*

$$\mathbb{E}_p[p(x)^k] = \sum_{j_1 + \dots + j_M = k} \binom{k}{j_1, \dots, j_M} \sum_{i=1}^M w_i \left( \frac{\mathcal{N}(0|\mu_i, \Sigma_i)}{\mathcal{N}(0|\mu, \Sigma)} \prod_{t=1}^M (w_t \mathcal{N}(0|\mu_t, \Sigma_t)^{j_t}) \right) \quad (27)$$

438 where  $\Sigma = (\Sigma_i^{-1} + \sum_{t=1}^M j_t \Sigma_t^{-1})^{-1}$  and  $\mu = \Sigma(\sum_{i=1}^M \mu_i + \sum_{t=1}^M j_t \Sigma_t^{-1} \mu_t)$ .

439 *Proof.* We will prove this statement by directly expanding out each term

$$\begin{aligned} \mathbb{E}_p[p(x)^k] &= \mathbb{E}_p \left[ \left( \sum_{t=1}^m w_t \mathcal{N}(x|\mu_t, \Sigma_t) \right)^k \right] \\ &= \mathbb{E}_p \left[ \sum_{j_1 + \dots + j_m = k} \binom{k}{j_1, \dots, j_m} \prod_{t=1}^m (w_t \mathcal{N}(x|\mu_t, \Sigma_t))^{j_t} \right] \\ &= \sum_{j_1 + \dots + j_m = k} \binom{k}{j_1, \dots, j_m} \prod_{t'=1}^m (w_{t'})^{j_{t'}} \sum_{i=1}^M w_i \int \mathcal{N}(x|\mu_i, \Sigma_i) \prod_{t=1}^m (\mathcal{N}(x|\mu_t, \Sigma_t))^{j_t} dx \end{aligned}$$

440 To combine the Gaussians under the integral, we appeal to the power of Gaussians (Lemma A.2.3)  
 441 and product of Gaussians (Lemma A.2.4)

$$\begin{aligned} &= \sum_{j_1 + \dots + j_m = k} \binom{k}{j_1, \dots, j_m} \prod_{t'=1}^m (w_{t'})^{j_{t'}} \sum_{i=1}^M w_i \int \mathcal{N}(x|\mu_i, \Sigma_i) \prod_{t=1}^m \frac{\mathcal{N}(x|\mu_t, \frac{1}{j_t} \Sigma_t)}{|j_t (2\pi \Sigma_t)^{j_t - 1}|^{1/2}} dx \\ &= \sum_{j_1 + \dots + j_m = k} \binom{k}{j_1, \dots, j_m} \prod_{t'=1}^m \frac{(w_{t'})^{j_{t'}}}{|j_t (2\pi \Sigma_t)^{j_t - 1}|^{1/2}} \sum_{i=1}^M w_i \int \frac{\mathcal{N}(0|\mu_i, \Sigma_i) \prod_{t=1}^m \mathcal{N}(0|\mu_t, \frac{1}{j_t} \Sigma_t)}{\mathcal{N}(0|\mu, \Sigma)} \mathcal{N}(x|\mu, \Sigma) dx \\ &= \sum_{j_1 + \dots + j_m = k} \binom{k}{j_1, \dots, j_m} \sum_{i=1}^M w_i \left( \frac{\mathcal{N}(0|\mu_i, \Sigma_i)}{\mathcal{N}(0|\mu, \Sigma)} \prod_{t=1}^m (w_t \mathcal{N}(0|\mu_t, \Sigma_t))^{j_t} \right) \end{aligned}$$

442 where  $\mu = \Sigma(\sum_{i=1}^M \mu_i + \sum_{t=1}^M j_t \Sigma_t^{-1} \mu_t)$  as defined from Lemma A.2.4. We see that we are left with  
 443 no integral and a closed form of the expectation of the powers of the GMM.  $\square$

444 **Theorem 4.2** (Convergence of  $\hat{H}_{N,a}^T(p(x))$ ). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a GMM and*  
 445 *choose a Taylor center such that  $a > \frac{1}{2} \max(p(x))$ . Then, for  $\hat{H}_{N,a}^T(p(x))$  defined in Eqn. (12)*

$$\lim_{N \rightarrow \infty} \hat{H}_{N,a}^T(p(x)) = H(p(x)) \quad (28)$$

446 *Proof.* We start out with the definition of entropy and introduce in the approximation discussed in  
 447 Lemma 3.2

$$\begin{aligned} H(p(x)) &= - \int \sum_{i=1}^M w_i q_i(x) \log(p(x)) dx = - \sum_{i=1}^M w_i \int q_i(x) \log(p(x)) dx \\ &= - \sum_{i=1}^M w_i \int q_i(x) \left( \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{n a^n} (p(x) - a)^n \right) dx \\ &= - \sum_{i=1}^M w_i \left( \log(a) + \int q_i(x) \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{n a^n} (p(x) - a)^n dx \right) \\ &= - \sum_{i=1}^M w_i \left( \log(a) + \int \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{n a^n} q_i(x) (p(x) - a)^n dx \right) \end{aligned}$$

448 We now wish to swap the order of integration and of the infinite summation as shown in Lemma A.2.1

$$\begin{aligned}
&= - \sum_{i=1}^M w_i \left( \log(a) + \sum_{i=1}^{\infty} \int \frac{(-1)^{n-1}}{na^n} q_i(x) (p(x) - a)^n dx \right) \\
&= - \sum_{i=1}^M w_i \left( \log(a) + \sum_{i=1}^{\infty} \frac{(-1)^{n-1}}{na^n} \int q_i(x) (p(x) - a)^n dx \right) \\
&= - \sum_{i=1}^M w_i \left( \log(a) + \sum_{i=1}^{\infty} \frac{(-1)^{n-1}}{na^n} \int q_i(x) \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} (p(x))^k dx \right) \\
&= - \sum_{i=1}^M w_i \left( \log(a) + \sum_{i=1}^{\infty} \frac{(-1)^{n-1}}{na^n} \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} \mathbb{E}_{q_i(x)} [(p(x))^k] \right)
\end{aligned}$$

449 We can compute  $\mathbb{E}_{q_i(x)} [(p(x))^k]$  using Lemma 4.1. The above term is equality for the Entropy,  
450 simply truncating the infinite summation gives a convergent approximation.

$$\hat{H}_{N,a}^T(p(x)) = - \sum_{i=1}^M w_i \left( \log(a) + \sum_{i=1}^N \frac{(-1)^{n-1}}{na^n} \sum_{k=0}^n \binom{n}{k} (-a)^{n-k} \mathbb{E}_{q_i(x)} [(p(x))^k] \right) \quad (29)$$

451

□

452 **Theorem 4.3** (Taylor Series is Lower Bound of Entropy). Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  and  
453  $a > \max(p(x))$ . Then, for all finite  $N$ ,

$$\hat{H}_{N,a}^T(p(x)) \leq H(p(x)) \quad (30)$$

454 *Proof.* If  $a > \max(p(x))$ , then we have the following lower bound

$$\begin{aligned}
H(p(x)) &= - \int p(x) \log p(x) dx = - \int p(x) \left( \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{na^n} (p(x) - a)^n \right) dx \\
&= - \int p(x) \left( \log(a) + \sum_{n=1}^{\infty} \frac{(-1)^{n-1} (-1)^n}{na^n} (a - p(x))^n \right) dx \\
&= - \log(a) - \int p(x) \left( \sum_{n=1}^{\infty} \frac{-1}{na^n} (a - p(x))^n \right) dx \\
&= - \log(a) + \int p(x) \left( \sum_{n=1}^{\infty} \frac{1}{na^n} (a - p(x))^n \right) dx \\
&\geq - \log(a) + \int p(x) \left( \sum_{n=1}^N \frac{1}{na^n} (a - p(x))^n \right) dx = \hat{H}_N(p(x))
\end{aligned}$$

455 since every term in the summation is positive due to  $a > p(x) \forall x$ , then truncating the series only  
456 removes positive terms, leaving us with a lower bound. □

457 **Theorem 4.4** (Upper bound on maximum of a GMM). Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$ , then

$$\max(p(x)) \leq a = \sum_i^M w_i |2\pi \Sigma_i|^{-\frac{1}{2}} \quad (31)$$

458 *Proof.* We need to find an upper bound on  $\max(p(x))$

$$\begin{aligned}
\max(p(x)) &= \max \left( \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i) \right) \\
&\leq \sum_{i=1}^M w_i \max(\mathcal{N}(x|\mu_i, \Sigma_i)) = \sum_{i=1}^M w_i |2\pi \Sigma_i|^{-\frac{1}{2}}
\end{aligned}$$

459 We simply have bound the maximum of the combination by combining the maximum of every  
 460 component in the GMM.  $\square$

461 **Theorem 4.5** (Convergence of  $\hat{H}_{N,a}^L(p(x))$ ). *Let  $p(x) = \sum_{i=1}^M w_i \mathcal{N}(x|\mu_i, \Sigma_i)$  be a GMM and*  
 462 *choose an interval such that  $a > \max(p(x))$ . Then for  $\hat{H}_{N,a}^L(p(x))$  defined in Eqn. (17)*

$$\lim_{N \rightarrow \infty} \hat{H}_{N,a}^L(p(x)) = H(p(x)) \quad (32)$$

463 *Proof.* We start out with the definition of entropy and introduce in the approximation discussed in  
 464 Lemma 3.3

$$\begin{aligned} H(p(x)) &= - \int \sum_{i=1}^M w_i q_i(x) \log(p(x)) dx = - \sum_{i=1}^M w_i \int q_i(x) \log(p(x)) dx \\ &= - \sum_{i=1}^M w_i \int q_i(x) \sum_{i=0}^{\infty} \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) dx \\ &= - \sum_{i=1}^M w_i \sum_{i=0}^{\infty} \int q_i(x) \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) dx \end{aligned}$$

465 We swapped the order of integration and of the infinite summation as shown in Lemma A.2.2. We  
 466 simply computation that are recreated in Theorem 3.3

$$\begin{aligned} &= - \sum_{i=0}^{\infty} \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} \sum_{i=1}^M w_i \int q_i(x) L_{[0,a],n}(p(x)) dx \\ &= - \sum_{i=0}^{\infty} (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! ((j+1)!)^2} \sum_{i=1}^M w_i \mathbb{E}_{q_i(x)} \left[ \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} p(x)^k \right] \\ &= - \sum_{i=0}^{\infty} (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! ((j+1)!)^2} \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} \sum_{i=1}^M w_i \mathbb{E}_{q_i(x)} [p(x)^k] \end{aligned}$$

467 We can compute  $\mathbb{E}_{q_i(x)} [p(x)^k]$  using Lemma 4.1. Simply truncating the infinite summation gives  
 468 the approximation.

$$\hat{H}_{N,a}^L(p(x)) = - \sum_{i=0}^N (2n+1) \sum_{j=0}^n \frac{(-1)^{n+j} (n+j)! ((j+1) \log(a) - 1)}{(n-j)! ((j+1)!)^2} \sum_{k=0}^n \frac{(-1)^{n+k} (n+k)!}{(n-k)! (k!)^2 a^k} \sum_{i=1}^M w_i \mathbb{E}_{q_i(x)} [p(x)^k]$$

469 which  $\lim_{N \rightarrow \infty} \hat{H}_{N,a}^L(p(x)) = H(p(x))$  as the above series is exactly equal to the entropy.  $\square$

470 Here, we address a few of the assumptions we made in the above derivations. We start with the ability  
 471 to swap the order of the integral and infinite sum for the Taylor series.

472 **Lemma A.2.1** (Swapping Integral and Infinite Sum (Taylor)). *Let  $a > \frac{1}{2} \max(p(x))$ , then*

$$\int \sum_{i=1}^{\infty} \frac{(-1)^{n-1}}{n a^n} q_i(x) (p(x) - a)^n dx = \sum_{i=1}^{\infty} \int \frac{(-1)^{n-1}}{n a^n} q_i(x) (p(x) - a)^n dx$$

473 *Proof.* For simplicity of notation, let  $c = \sup \left| \frac{p(x)-a}{a} \right| < 1$  since  $a > \frac{1}{2} \max(p(x))$ . We then appeal  
 474 to Fubini-Tonelli theorem which states that if  $\int \sum |f_n(x)| dx < \infty$  or if  $\sum \int |f_n(x)| dx < \infty$ , then  
 475  $\int \sum f_n(x) dx = \sum \int f_n(x) dx$ .

$$\begin{aligned} \sum_{i=1}^{\infty} \int \left| \frac{(-1)^{n-1}}{n a^n} q_i(x) (p(x) - a)^n \right| dx &= \sum_{i=1}^{\infty} \frac{1}{n} \int q_i(x) \left( \frac{|p(x) - a|}{a} \right)^n dx \\ &\leq \sum_{i=1}^{\infty} \frac{1}{n} \int q_i(x) (c)^n dx \\ &= \sum_{i=1}^{\infty} \frac{c^n}{n} \int q_i(x) dx = \sum_{i=1}^{\infty} \frac{c^n}{n} < \infty \end{aligned}$$

476 We know that  $\sum_{i=1}^{\infty} \frac{c^n}{n} < \infty$  because of the ratio test again

$$\lim_{n \rightarrow \infty} \left| \frac{c^{n+1}}{n+1} \frac{n}{c^n} \right| = \lim_{n \rightarrow \infty} \frac{n}{n+1} c = c < 1$$

477 So we see that Fubini-Tonelli holds so

$$\int \sum_{i=1}^{\infty} \frac{(-1)^{n-1}}{na^n} q_i(x) (p(x) - a)^n dx = \sum_{i=1}^{\infty} \int \frac{(-1)^{n-1}}{na^n} q_i(x) (p(x) - a)^n dx$$

478 □

479 We now consider the case for the Legendre series

480 **Lemma A.2.2** (Swapping Integral and Infinite Sum (Legendre)). *Let  $a > \max(p(x))$ , then*

$$\int q_i(x) \sum_{i=1}^{\infty} \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) dx = \sum_{i=1}^{\infty} \int q_i(x) \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) dx$$

481 *Proof.* We again appeal to Fubini-Tonelli. We will use that  $\left| \frac{L_{[0,a],n}(p(x))}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} \right| \leq 1$  as it is the  
482 orthonormal polynomials and then use Cauchy-Schwartz on the remaining term

$$\begin{aligned} & \sum_{i=1}^{\infty} \int \left| q_i(x) \frac{\langle \log(y), L_{[0,a],n}(y) \rangle}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} L_{[0,a],n}(p(x)) \right| dx \\ &= \sum_{i=1}^{\infty} \langle \log(y), L_{[0,a],n}(y) \rangle \int \left| \frac{q_i(x) L_{[0,a],n}(p(x))}{\langle L_{[0,a],n}(y), L_{[0,a],n}(y) \rangle} \right| dx \\ &\leq \sum_{i=1}^{\infty} \|\log(y)\|^2 \|L_{[0,a],n}(y)\|^2 \int q_i(x) dx \\ &= \sum_{i=1}^{\infty} a((\log(a) - 2) \log(a) + 2) \left( \frac{a}{2n+1} \right)^2 \int N(x|\mu_i, \Sigma_i) dx \\ &= a^3((\log(a) - 2) \log(a) + 2) \sum_{i=1}^{\infty} \frac{1}{(2n+1)^2} < \infty \end{aligned}$$

483 So we see that since the absolute value is finite, then Fubini-Tonelli applies and we can swap the  
484 order of the integral and infinite summation. □

485 The next thing we show is the relations we used for powers of Gaussians.

486 **Lemma A.2.3.** *Powers of a Gaussian*

487

$$\mathcal{N}(x|\mu, \Sigma)^n = |n(2\pi\Sigma)^{n-1}|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n}\Sigma\right)$$

488 *Proof.* We are going to do an inductive proof and rely on the well known relation of products of  
489 Gaussians

$$\mathcal{N}(x|a, A)\mathcal{N}(x|b, B) = \mathcal{N}(x|d, D)\mathcal{N}(a|b, A+B)$$

490 Where  $D = (A^{-1} + B^{-1})^{-1}$  and  $d = D(A^{-1}a + B^{-1}b)$ .

491 **Base Case:**  $n = 2$

$$\mathcal{N}(x|\mu, \Sigma)^2 = \mathcal{N}(x|\mu, \Sigma)\mathcal{N}(x|\mu, \Sigma) \tag{33}$$

$$= \mathcal{N}(x|\mu, \frac{1}{2}\Sigma)\mathcal{N}(\mu|\mu, 2\Sigma) = |2\pi\Sigma|^{-\frac{1}{2}} \mathcal{N}(x|\mu, \frac{1}{2}\Sigma) \tag{34}$$

492 In this case, we get that  $D = (\Sigma^{-1} + \Sigma^{-1})^{-1} = \frac{1}{2}\Sigma$  and  $d = \frac{1}{2}\Sigma(\Sigma^{-1}\mu + \Sigma^{-1}\mu) = \mu$ . We also  
493 see that  $\mathcal{N}(\mu|\mu, 2\Sigma)$  is being evaluated at its maximum, which just leaves the scaling term out front

494 of the exponential in the Gaussian,  $|2\pi\Sigma|^{-\frac{1}{2}}$ .

495 **Inductive step:** Assume that  $N(x|\mu, \Sigma)^n = |n(2\pi\Sigma)^{n-1}|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n}\Sigma\right)$ , then we wish to show  
 496 that  $N(x|\mu, \Sigma)^{n+1} = |(n+1)(2\pi\Sigma)^n|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n+1}\Sigma\right)$

$$\begin{aligned} N(x|\mu, \Sigma)^{n+1} &= N(x|\mu, \Sigma)^n N(x|\mu, \Sigma) \\ &= |n(2\pi\Sigma)^{n-1}|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n}\Sigma\right) N(x|\mu, \Sigma) \\ &= |n(2\pi\Sigma)^{n-1}|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n+1}\Sigma\right) N\left(\mu|\mu, \frac{n+1}{n}\Sigma\right) \\ &= |n(2\pi\Sigma)^{n-1}|^{-\frac{1}{2}} \left|2\pi\frac{n+1}{n}\Sigma\right|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n+1}\Sigma\right) \\ &= |(n+1)(2\pi\Sigma)^n|^{-\frac{1}{2}} \mathcal{N}\left(x|\mu, \frac{1}{n+1}\Sigma\right) \end{aligned}$$

497 Here,  $D = \left(\left(\frac{1}{n}\Sigma\right)^{-1} + \Sigma^{-1}\right)^{-1} = \frac{1}{n+1}\Sigma$  and  $d = \frac{1}{n+1}\Sigma\left(\left(\frac{1}{n}\Sigma\right)^{-1}\mu + \Sigma^{-1}\mu\right) = \mu$ .  $\square$

498 We finally show the product of Gaussians that we used. We keep the exact same notation used in the  
 499 derivation of the entropy Taylor series so the terms may be more easily identified.

500 **Lemma A.2.4.** *Product of a Gaussians*

501

$$\mathcal{N}(x|\mu_i, \Sigma_i) \prod_{t=1}^m \mathcal{N}\left(x|\mu_t, \frac{1}{j_t}\Sigma_t\right) = \mathcal{N}(0|\mu_i, \Sigma_i) \prod_{t=1}^m \mathcal{N}\left(0 + \mu_t, \frac{1}{j_t}\Sigma_t\right) \frac{N(x|\mu, \Sigma)}{N(0|\mu, \Sigma)}$$

502 where  $\Sigma = (\Sigma^{-1} + \sum_{t=1}^m j_t \Sigma_t^{-1})^{-1}$  and  $\mu = \Sigma (\Sigma_i^{-1} \mu_i + \sum_{t=1}^m j_t \Sigma_t^{-1} \mu_t)$

503 *Proof.* We will simply expand out the product of Gaussians, collect like terms, complete the square, and then  
 504 recollect exponentials into Gaussians evaluated at 0.

$$\begin{aligned} \mathcal{N}(x|\mu_i, \Sigma_i) \prod_{t=1}^m \mathcal{N}\left(x|\mu_t, \frac{1}{j_t}\Sigma_t\right) &= |2\pi\Sigma_i|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}(x - \mu_i)^T \Sigma_i^{-1} (x - \mu_i)\right\} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}(x - \mu_t)^T j_t \Sigma_t^{-1} (x - \mu_t)\right\} \\ &= |2\pi\Sigma_i|^{-\frac{1}{2}} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(x^T \Sigma_i^{-1} x - 2x^T \Sigma_i^{-1} \mu_i + \mu_i^T \Sigma_i^{-1} \mu_i + \sum_{t=1}^m x^T j_t \Sigma_t^{-1} x - 2x^T j_t \Sigma_t^{-1} \mu_t + \mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} \\ &= |2\pi\Sigma_i|^{-\frac{1}{2}} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(x^T \left(\Sigma_i^{-1} + \sum_{t=1}^m j_t \Sigma_t^{-1}\right) x - 2x^T \left(\Sigma_i^{-1} \mu_i + \sum_{t=1}^m j_t \Sigma_t^{-1} \mu_t\right) + \mu_i^T \Sigma_i^{-1} \mu_i + \sum_{t=1}^m \mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} \end{aligned}$$

505 Now we let  $\Sigma = (\Sigma^{-1} + \sum_{t=1}^m j_t \Sigma_t^{-1})^{-1}$  and  $\mu = \Sigma (\Sigma_i^{-1} \mu_i + \sum_{t=1}^m j_t \Sigma_t^{-1} \mu_t)$

$$\begin{aligned} &= |2\pi\Sigma_i|^{-\frac{1}{2}} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(x^T \Sigma^{-1} x - 2x^T \Sigma^{-1} \mu + \mu_i^T \Sigma_i^{-1} \mu_i + \sum_{t=1}^m \mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} \\ &= |2\pi\Sigma_i|^{-\frac{1}{2}} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(x^T \Sigma^{-1} x - 2x^T \Sigma^{-1} \mu + \mu^T \Sigma \mu - \mu^T \Sigma \mu + \mu_i^T \Sigma_i^{-1} \mu_i + \sum_{t=1}^m \mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} \\ &= |2\pi\Sigma_i|^{-\frac{1}{2}} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \frac{|2\pi\Sigma|^{1/2}}{|2\pi\Sigma|^{1/2}} \exp\left\{-\frac{1}{2}\left((x - \mu)^T \Sigma^{-1} (x - \mu)\right)\right\} \exp\left\{-\frac{1}{2}\left(-\mu^T \Sigma \mu + \mu_i^T \Sigma_i^{-1} \mu_i + \sum_{t=1}^m \mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} \\ &= |2\pi\Sigma_i|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(\mu_i^T \Sigma_i^{-1} \mu_i\right)\right\} \prod_{t=1}^m \left|\frac{2\pi}{j_t}\Sigma_t\right|^{-\frac{1}{2}} \exp\left\{-\frac{1}{2}\left(\mu_t^T j_t \Sigma_t^{-1} \mu_t\right)\right\} |2\pi\Sigma|^{1/2} \exp\left\{-\frac{1}{2}\left(-\mu^T \Sigma \mu\right)\right\} \mathcal{N}(x|\mu, \Sigma) \\ &= \mathcal{N}(0|\mu_i, \Sigma_i) \prod_{t=1}^m \mathcal{N}\left(0|\mu_t, \frac{1}{j_t}\Sigma_t\right) \frac{N(x|\mu, \Sigma)}{N(0|\mu, \Sigma)} \end{aligned}$$

506

$\square$

507 **A.4 Extensions to Cross-Entropy**

508 Our main results on GMM entropy approximation also extend to the cross-entropy between  
 509 different GMMs. We can instead consider  $H_p(q(x)) = -\mathbb{E}_{p(x)}[\log(q(x))]$  where  $p(x) =$   
 510  $\sum_{i=1}^{\widehat{M}} \widehat{w}_i \mathcal{N}(x|\widehat{\mu}_i, \widehat{\Sigma}_i)$  and  $q(x) = \sum_{i=1}^{\widehat{M}} \widehat{w}_i \mathcal{N}(x|\widehat{\mu}_i, \widehat{\Sigma}_i)$ . The series representations of  $\log(\cdot)$  stay  
 511 unchanged however we must choose a center point that allows the series to converge with respect  
 512 to the inner GMM,  $q(x)$ . This means Theorem 3.2, Theorem 3.3, Theorem 4.2, Theorem 4.3, and  
 513 Theorem 4.5 need to be reformulated so that  $a > \max(q(x))$  (or the respective  $a > \frac{1}{2} \max(q(x))$ ). We  
 514 do not formally restate these theorems here for brevity. We instead provide a sketch of how the results  
 515 extend to the cross-entropy setting. Choosing the bounding  $a$  of the max found in Theorem 4.4 can  
 516 simply be altered so that

$$\max(q(x)) \leq a = \sum_i^{\widehat{M}} \widehat{w}_i \left| 2\pi \widehat{\Sigma}_i \right|^{-\frac{1}{2}}$$

517 The analogous proofs will all hold in this case. The final alteration that needs to be made is to  
 518 Theorem 4.1. Again, following the exact same proof, just switching notation and being careful, one  
 519 can find that

$$\mathbb{E}_p[q(x)^k] = \sum_{j_1 + \dots + j_{\widehat{M}} = k} \binom{k}{j_1, \dots, j_{\widehat{M}}} \sum_{i=1}^{\widehat{M}} \widehat{w}_i \left( \frac{\mathcal{N}(0|\widehat{\mu}_i, \widehat{\Sigma}_i)}{\mathcal{N}(0|\mu, \Sigma)} \prod_{t=1}^{\widehat{M}} (\widehat{w}_t \mathcal{N}(0|\widehat{\mu}_t, \widehat{\Sigma}_t)^{j_t}) \right) \quad (35)$$

520 where  $\Sigma = (\widehat{\Sigma}_i^{-1} + \sum_{t=1}^{\widehat{M}} j_t \widehat{\Sigma}_t^{-1})^{-1}$  and  $\mu = \Sigma(\widehat{\Sigma}_i^{-1} \widehat{\mu}_i + \sum_{t=1}^{\widehat{M}} j_t \widehat{\Sigma}_t^{-1} \widehat{\mu}_t)$ . The result is the  
 521 same form, however has much more convoluted notation and hence dropped from the main paper as  
 522 an attempt to bring clarity to the methods being discuss without unnecessary notation.

523 **A.5 Discussion of Taylor Limit**

524 Computing the series in Eqn. (12) for higher orders can be computationally prohibitive. In particular,  
 525 the sum  $\sum_{j_1 + \dots + j_M = n}$  is over  $M$  integers summing to  $n$ , which is  $\mathcal{O}((n + M - 1)!)$ . In this section,  
 526 we provide an approximation that avoids explicit computation of this sum for higher orders. The  
 527 method is based on a polynomial fit of the convergence rate for the lower bound property discussed  
 528 in Theorem 4.3. We model this convergence as,

$$\hat{H}_n(p(x)) = \beta \cdot \alpha^n + \eta \quad (36)$$

529 where  $\hat{H}_n(p(x))$  is the  $n^{\text{th}}$  order Taylor approximation and we have dropped explicit dependence on  
 530 the center and notation of Taylor method for brevity. Further,  $\beta < 0$  and  $0 < \alpha < 1$  are convergence  
 531 constants and  $\eta$  is the estimated limiting value of the approximation. We require three consecutive  
 532 orders of our Taylor series approximation,  $\hat{H}_n(p(x))$ ,  $\hat{H}_{n+1}(p(x))$ , and  $\hat{H}_{n+2}(p(x))$ , to solve for the  
 533 three unknown parameters,

$$\begin{aligned} \hat{H}_n(p(x)) &= \beta \cdot \alpha^n + \eta \\ \hat{H}_{n+1}(p(x)) &= \beta \cdot \alpha^{n+1} + \eta \\ \hat{H}_{n+2}(p(x)) &= \beta \cdot \alpha^{n+2} + \eta. \end{aligned}$$

534 Given three equations and three unknowns we can solve for the approximate limiting entropy as,

$$\eta = \hat{H}_n - \frac{(\hat{H}_{n+1} - \hat{H}_n)^2}{\hat{H}_{n+2} - 2\hat{H}_{n+1} + \hat{H}_n} \quad (37)$$

535 This approach assumes that the Taylor series converges according to Eqn. (36), which is not the  
 536 case in general. Identifying the exact rate of convergence is a topic of future work. But this simple  
 537 approximation has shown higher accuracy in practice with negligible additional computation, as  
 538 shown in the experiments of Sec. 6. With slight abuse of terminology we refer to this approach as the  
 539 *Taylor limit*. We do not apply this method to the Legendre approximation (Eqn. (17)) as it doesn't  
 540 maintain a lower bound during its convergence. Equivalent methods have been considered to model  
 541 the potential oscillation convergence however in practice, we do not find an increase in accuracy.

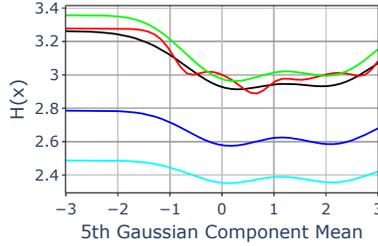


Figure 6: **Well-Behaved GMM** example is plotted on the left. The entropy of a five-component, bivariate GMM is plotted as a function of the location of the fifth component  $\mu_5 \in [-3, 3]$ . We show the true entropy, the 3<sup>rd</sup> order of Huber et al.’s approximation and our two methods, along with our approximate limit.

## 542 A.6 Recreation of Huber et al Experiment

543 Here we reproduce the experiment of [15] of a five-component bivariate GMM with uniform weights  
 544  $w_i = 0.2$  for  $i = 1, \dots, 5$ ,  $\mu_1 = [0, 0]^T$ ,  $\mu_2 = [3, 2]^T$ ,  $\mu_3 = [1, 0.5]^T$ ,  $\mu_4 = [2.5, 1.5]^T$ ,  
 545  $\mu_5 = c \cdot [1, 1]^T$ ,  $\Sigma_1 = \text{diag}(0.16, 1)$ ,  $\Sigma_2 = \text{diag}(1, 0.16)$ , and  $\Sigma_3 = \Sigma_4 = \Sigma_5 = \text{diag}(0.5, 0.5)$ .  
 546 We vary the position of the fifth mean ( $\mu_5$ ) in the range  $[-3, 3]$ . Fig. 6 (left) reports the third order  
 547 Taylor approximations from both Taylor approximations, the Legendre approximation, as well as our  
 548 approximate limiting method.

549 Huber et al. is accurate in the well-behaved case, but does not have any convergence guarantees nor  
 550 is it a bound. Our proposed Taylor approximation sacrifices some accuracy, but is always a lower  
 551 bound (Theorem 4.3) and is convergent (Theorem 4.2). We also note that our naïve limit method does  
 552 gain us substantial accuracy and is still a lower bound—though we have not proven the bound property  
 553 for this approximation. We notice that our Legendre approximation has comparable accuracy to  
 554 that of Huber et al. in this well behaved case but has the advantage that it is guaranteed to converge  
 555 (Theorem 4.5) and that we can compute higher order approximations that are difficult to define for  
 556 the Huber et al. approximation.

## 557 A.7 Nonparametric Variational Inference

558 Consider a target density  $p(x, \mathcal{D})$  with latent variables  $x$  and observations  $\mathcal{D}$ . The NPV approach [11]  
 559 optimizes the evidence lower bound (ELBO),

$$\log p(x) \geq \max_q H_q(p(x, \mathcal{D})) - H_q(q(x)) \equiv \mathcal{L}(q) \quad (38)$$

560 w.r.t. a  $m$ -component GMM variational distribution  $q(x) = \frac{1}{N} \sum_{i=1}^m \mathcal{N}(x | \mu_i, \sigma_i^2 I_d)$ . The GMM  
 561 entropy lacks a closed-form so NPV applies Jensen’s lower bound as an approximation,

$$H_q(q(x)) = -\mathbb{E}_q [\log(q(x))] \geq -\sum_{i=1}^M w_i \log(\mathbb{E}_{q_i} [q(x)]) = \hat{H}_q^J(q(x)) \quad (39)$$

562 The cross entropy also lacks a closed-form, so NPV approximates this term using the analogous  
 563 Huber et al. Taylor approximation. Specifically, NPV expands the log density around the means of  
 564 each GMM component as,

$$H_q(p(x)) \approx -\sum_{i=1}^M w_i \sum_{n=0}^N \frac{\nabla^2 \log(p(\mu_i))}{n!} \mathbb{E}_{q_i} [(x - \mu_i)^n] = \hat{H}_{N,q}^H(p(x)) \quad (40)$$

565 However, Eqn. (19) is subject to the divergence criterion of Theorem 3.1 if  $2p(\mu_i) \leq \max(p(x))$ .  
 566 This approximation is often known as the multivariate delta method for moments. The authors use  
 567 these approximations of the entropy and cross entropy to create the following two approximation of  
 568 the ELBO.

$$\mathcal{L}_1(q) = \hat{H}_{1,q}^H(p(x)) - \hat{H}_q^J(q(x)) \quad \mathcal{L}_2(q) = \hat{H}_{2,q}^H(p(x)) - \hat{H}_q^J(q(x)) \quad (41)$$

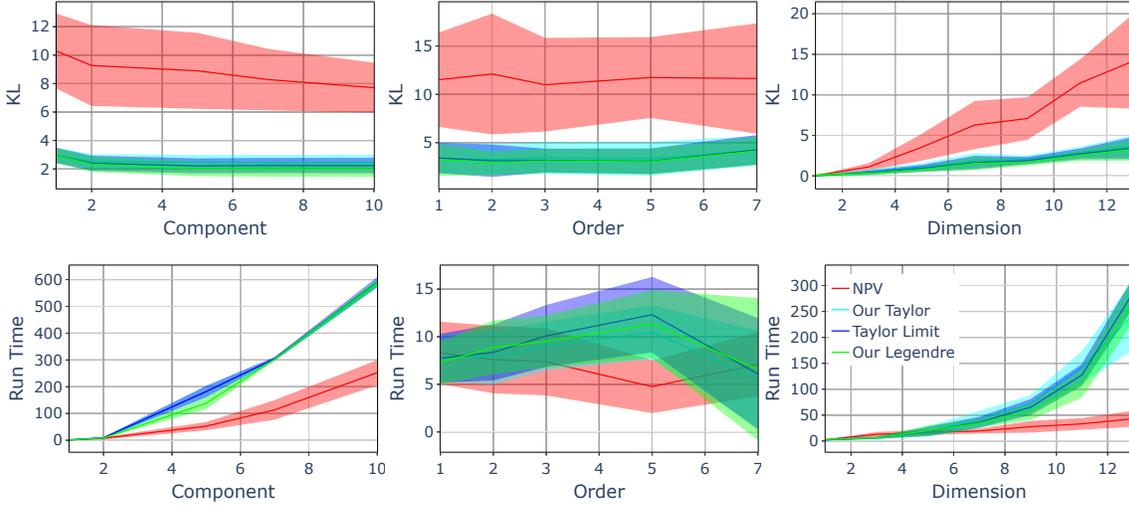


Figure 7: The above figures show the accuracy and computation time of each method across varying components, orders, and dimensions in approximating a multivariate mixture T distribution with a GMM. Our method consistently improves accuracy significantly. Higher order approximations do not increase computation time (bottom middle), and even low order approximations provide substantial improvement (top middle). Most accuracy improvements are achieved with a small number of component, unlike NPV (top left) which continues to need higher number of components to see good accuracy return. Our computation time increases with dimension due to less salable cross-entropy approximation with Gauss-Hermite quadrature (bottom right) however the guaranteed convergence of the approximation seems to have a drastic improvement on accuracy (top right).

569 Gershman et al. (2012) use the two approximations because optimizing  $\hat{H}_{2,q}^H(p(x))$  with respect to  
570 the mean components,  $\mu_i$  requires computing a third order, multivariate derivative of  $\log(p(x))$  which  
571 is computationally expensive. The authors iterate between optimizing the mean components,  $\mu_i$ ,  
572 using  $\mathcal{L}_1(q)$  and optimizing the variance components,  $\sigma_i^2$ , using  $\mathcal{L}_2(q)$  until the second approximate  
573 appropriately converges  $\delta\mathcal{L}_2(q) < .0001$ .

574 **In our approach**, we will highlight and address three problems with the NPV algorithm; the potential  
575 divergence of  $\hat{H}_{N,q}^H(p(x))$ , the inconsistent ELBO approximations, and the poor estimation of the  
576 GMM entropy via  $\hat{H}_q^J(q(x))$ . To address the potential divergence of  $\hat{H}_{N,q}^H(p(x))$ , we will take  
577 motivation from the results found in [23] and use a 2 point Gauss-Hermite quadrature method  
578 to approximate  $H_q(p(x))$ . This method will be a limiting factor in scaling the NPV algorithm  
579 in dimension, however it guarantees that the cross-entropy approximation will not diverge. This  
580 alteration leads to a solution for the inconsistency of the ELBO approximations. Since the quadrature  
581 method does not require any derivatives of  $\log(p(x))$  w.r.t. the mean components of the GMM,  
582 we can now optimize the means and variances simultaneously on the same ELBO approximation.  
583 Finally, Jensen’s inequality is a very poor approximation for entropy in general, instead we will use  
584 the three methods we have introduced, Taylor, Taylor limit, and Legendre, as the GMM entropy  
585 approximations for higher accuracy. Fig. 4 shows an approximation of a two dimensional, three  
586 component mixture Student T distribution using a five component GMM in the traditional NPV,  
587 using our Taylor approximation and our Legendre approximation.

588 **The results**, as seen in Fig. 5, highlight the accuracy and computation time of each method versus the  
589 number of components used, the order of our polynomial approximation used, and the dimension of  
590 the GMM. The accuracy is the same as seen in Section 6, the new information here is the computation  
591 time of each method. We see that the order of the method has very little impact on the computation  
592 time of our algorithm. We even see most of the accuracy improvement at around order 2 or 3 so  
593 staying in a low order approximation seems advisable. We do see that the component does increase  
594 our time by a bit compared to that of traditional NPV. The source of the computation time increase  
595 in our methods comes from more iterations in the optimization. Each evaluation of the (ELBO)  
596 approximator are near equivalent but since we are converging to a better optimum, this take more

597 iteration steps than NPV. Finally, we see that dimension does have a large impact on our method.  
598 The source of this computation increase come from our Gauss-Hermite approximation of the cross  
599 entropy. The number of quadrature points used 2 per dimension  $D$ , so we are computing with  $2^D$   
600 points, which clearly scale poorly with dimension. We are seeking better ways of computing the  
601 cross-entropy of a GMM with any non-GMM function that is both convergent and computationally  
602 fast, however this was not the focus of the paper and no method was considered yet.