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# Which graphical models are difficult to learn?

## Supplementary material

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In this document we report some of the calculations and proofs omitted from the main paper. References to equations and theorems in the main paper are in boldface.

### 1 Thresholding algorithm

In the following we let  $C_{ij} \equiv \mathbb{E}_{G,\theta}\{x_i x_j\}$  where expectation is taken with respect to the Ising model **(1)**.

Before proving Theorem **1.1** we start with an easy related lemma.

**Lemma 1.0.** If  $G$  is a tree, and  $\tau(\theta) = (\tanh \theta + \tanh^2 \theta)/2$ , then

$$n_{\text{Thr}(\tau)}(G, \theta) \leq \frac{8}{(\tanh \theta - \tanh^2 \theta)^2} \log \frac{2p}{\delta}. \quad (1)$$

*Proof.* (Lemma 1.0) If  $G$  is a tree then  $C_{ij} = \tanh \theta$  for all  $(ij) \in E$  and  $C_{ij} \leq \tanh^2 \theta$  for all  $(ij) \notin E$ . The probability that  $\text{Thr}(\tau)$  fails is

$$1 - P_{\text{succ}} = \mathbb{P}_{n,G,\theta}\{\hat{C}_{ij} < \tau \text{ for some } (i,j) \in E \text{ or } \hat{C}_{ij} \geq \tau \text{ for some } (i,j) \notin E\}. \quad (2)$$

Let  $\tau = (\tanh \theta + \tanh^2 \theta)/2$ . Applying Azuma-Hoeffding inequality to  $\hat{C}_{ij}$  followed by union bound over the edges, we bound this probability by

$$P_{\text{succ}} \geq 1 - p^2 e^{-\frac{1}{8}n(\tanh \theta - \tanh^2 \theta)^2}. \quad (3)$$

Imposing the right hand side to be larger than  $\delta$  proves our result.  $\square$

*Proof.* (Theorem **1.1**) We will prove that, for  $\theta < \text{arctanh}(1/(2\Delta))$ ,  $C_{ij} \geq \tanh \theta$  for all  $(i,j) \in E$  and  $C_{ij} \leq 1/(2\Delta)$  for all  $(ij) \notin E$ . In particular  $C_{ij} < C_{kl}$  for all  $(i,j) \notin E$  and all  $(k,l) \in E$ . The theorem follows from this fact via union bound and Azuma-Hoeffding inequality as in the proof of Theorem 1.0.

The bound  $C_{ij} \geq \tanh \theta$  for  $(ij) \in E$  is a direct consequence of Griffiths inequality: compare the expectation of  $x_i x_j$  in  $G$  with the same expectation in the graph that only includes edge  $(i,j)$ .

The second bound is derived using the technique of **[16]**, i.e., bound  $C_{ij}$  by the generating function for self-avoiding walks on the graphs from  $i$  to  $j$ . More precisely, assume  $l = \text{dist}(i,j)$  and denote by  $N_{ij}(k)$  the number of self avoiding walks of length  $k$  between  $i$  and  $j$  on  $G$ . Then **[16]** proves that

$$C_{ij} \leq \sum_{k=l}^{\infty} (\tanh \theta)^k N_{ij}(k) \leq \sum_{n=l}^{\infty} \Delta^{k-1} (\tanh \theta)^k \leq \frac{\Delta^{l-1} (\tanh \theta)^l}{1 - \Delta \tanh \theta} \leq \frac{\Delta (\tanh \theta)^2}{1 - \Delta \tanh \theta} \quad (4)$$

If  $\theta < \text{arctanh}(1/(2\Delta))$  the above implies  $C_{ij} \leq 1/(2\Delta)$  which is our claim.  $\square$

*Proof.* (Theorem 1.2) The theorem is proved by constructing  $G$  as follows: sample a uniformly random regular graph of degree  $\Delta$  over the  $p - 2$  vertices  $\{1, 2, \dots, p - 2\} \equiv [p - 2]$ . Add an extra edge between nodes  $p - 1$  and  $p$ . The resulting graph is not connected. We claim that for  $\theta > K/\Delta$  and with probability converging to 1 as  $p \rightarrow \infty$ , there exist  $i, j \in [p - 2]$  such that  $(i, j) \notin E$  and  $C_{ij} > C_{p-1,p}$ . As a consequence, thresholding fails.

Obviously  $C_{p-1,p} = \tanh \theta$ . Choose  $i \in [p - 2]$  uniformly at random, and  $j$  a node at a fixed distance  $t$  from  $i$ . We can compute  $C_{ij}$  as  $p \rightarrow \infty$  using the same local weak convergence result as in the proof of Lemma 3.3. Namely,  $C_{ij}$  converges to the correlation between the root and a leaf node in the tree Ising model (16). In particular one can show that

$$\lim_{p \rightarrow \infty} C_{ij} \geq m(\theta)^2, \quad (5)$$

where  $m(\theta) = \tanh(\Delta h^*/(\Delta - 1))$  and  $h^*$  is the unique positive solution of  $h = (\Delta - 1) \operatorname{atanh}\{\tanh \theta \tanh h\}$ .

The proof is completed by showing that  $\tanh \theta < m(\theta)^2$  for all  $\theta > K/\Delta$ .  $\square$

## 2 Regularized logistic regression

*Proof.* (Lemma 3.1) We outline here the upper bound on the term  $R^n$ .

Since  $\hat{\theta}_{SC} = 0$  an application of the mean value theorem yields  $|[R^n]_j| \leq 2\Delta \|\hat{\theta}_S - \theta_S^*\|_2^2$ . Now  $\sigma_{\min}(Q^*) \leq 1$  so the event  $\mathcal{E}$  guarantees that  $\sigma_{\min}(Q^{n*}) \leq 2$ . Using Lemma 3 from [7] we can write

$$\|\hat{\theta}_S - \theta_S^*\|_2 \leq \frac{1}{\Delta^{3/2}} \left( 1 - \sqrt{1 - \lambda \frac{8\Delta^2}{C_{\min}} (1 + \|\frac{W_S^n}{\lambda}\|_{\infty})} \right). \quad (6)$$

If  $\mathcal{E}$  holds we can assume without loss of generality  $\|\frac{W_S^n}{\lambda}\|_{\infty} < 1$  and since  $1 - \sqrt{1 - x} \leq x, x \in [0, 1]$  the theorem's assumption on  $\lambda$  makes both  $\frac{\Delta}{C_{\min}} \|\frac{R_S^n}{\lambda}\|_{\infty}$  and  $|\frac{R_v^n}{\lambda}|$  smaller than  $\epsilon/8$ .  $\square$

*Proof.* (Lemma 3.3) We outline here some of the calculations with respect to the tree model (16). An important property that follows from the fixed point equation  $h = (\Delta - 1) \operatorname{atanh}\{\tanh \theta \tanh h\}$  is that, if  $g(\underline{x}_{T(t)})$  is a function of the variables in  $T(t)$  then

$$\mathbb{E}_{T(t), \theta, +}\{g(\underline{X}_{T(t)})\} = \mathbb{E}_{T(t+1), \theta, +}\{g(\underline{X}_{T(t)})\}, \quad (7)$$

with the obvious identification of  $T(t)$  as a subtree of  $T(t + 1)$ .

Let  $r$  be a uniformly random vertex in  $G$  and  $i \neq j$  two neighbors of  $r$ . Using the local weak convergence property (17) with  $t = 1$  we get

$$\lim_{p \rightarrow \infty} (Q_{SS}^*)_{ii} \equiv a = \mathbb{E}_{T(1), \theta, +}\left(\frac{1}{\cosh^2 \theta M}\right), \quad (8)$$

$$\lim_{p \rightarrow \infty} (Q_{SS}^*)_{ij} \equiv b = \mathbb{E}_{T(1), \theta, +}\left(\frac{X_i X_j}{\cosh^2 \theta M}\right), \quad (9)$$

where  $M \equiv \sum_{i \in \partial T(1)} X_i$  is the sum of the variables on the leaves of a depth 1 tree, and  $i, j \in \partial T(1)$ . Let  $c$  and  $d$  be defined by

$$\lim_{p \rightarrow \infty} (Q_{SS}^*)_{ii}^{-1} \equiv c, \quad (10)$$

$$\lim_{p \rightarrow \infty} (Q_{SS}^*)_{ij}^{-1} \equiv d. \quad (11)$$

Finally, for  $r'$  at distance  $t$  from  $r$ , consider the  $\Delta$ -dimensional vector in

$$\lim_{p \rightarrow \infty} (Q_{S^c S}^*)_{r'} = F_S(t). \quad (12)$$

It can be shown that

$$F_S(t) = \mathbb{E}_{T(1), \theta, +}(X_{r'}) \mathbb{E}_{T(1), \theta, +}\left(\frac{X_i}{\cosh^2 \theta M}\right) + o_t(1) = f + o_t(1), \quad (13)$$

where  $r'$  is the root of a first tree, and  $i \in \partial T(1)$  is a leaf of the second tree. In particular  $F_S(t)$  has, for large  $t$ , asymptotically equal entries.

The final result is obtained by computing the quantities  $a, b, c, d, f$ . □